

# What's on the menu this week? – an appetizer

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# Contents

- 1 Preparing the food
- 2 The cutlery
- 3 Serving the food
  - The starters – they taste differently
  - The starters – more cutlery
  - Important ingredients for a connoisseur
  - Developing more taste
  - Will there be a main dish?
- 4 The main dish: something to chew on!

## Why did you get invited for dinner?

- The overall motivation is to continue the workshops of 1995–2011 in Enschede, Nectiny, Hannover, Hajek and Domažlice in order to make **progress** on several intriguing conjectures.
- These highly interrelated conjectures involve line graphs, claw-free graphs, cubic graphs, snarks, and concepts like Hamilton cycles, Hamilton-connectedness, dominating closed trails (circuits), and dominating cycles.
- Perhaps there is a link to double cycle covers and nowhere zero flows, or other conjectures; there is a link to the P vs NP millennium problem.
- In order to introduce the workshop topics, I was asked to repeat some background, and present a **survey** on some of the conjectures and their relationships.
- More details can be found in the **survey paper** copied from Graphs and Combinatorics.

## Starting from scratch: the first two conjectures

The following two conjectures were tossed in the eighties.

Matthews & Sumner, 1984:

### Conjecture (MS-Conjecture)

*Every 4-connected **claw-free graph** is **hamiltonian**.*

Thomassen, 1986:

### Conjecture (T-Conjecture)

*Every 4-connected **line graph** is **hamiltonian**.*

Let me start by explaining the **terminology** to understand the above statements and their relationship.

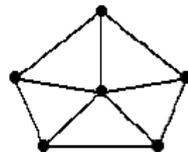
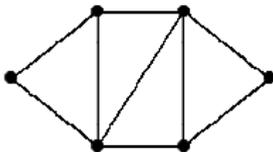
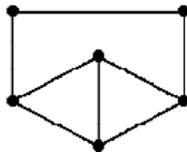
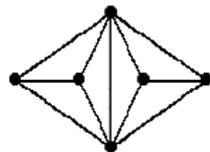
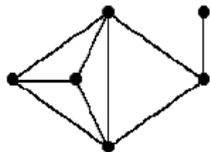
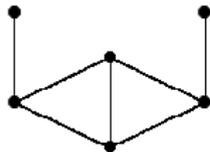
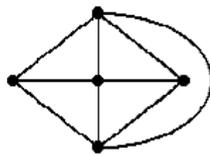
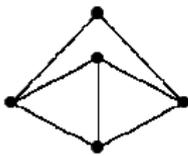
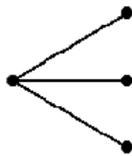
# Basic graph terminology

- All graphs in this talk are finite, undirected, loopless and the majority is simple (in some results we allow multiple edges) .
- We denote a graph  $G$  as  $G = (V, E)$ , where  $V = V(G)$  is the vertex set and  $E = E(G)$  is the edge set.
- A graph is **hamiltonian** if it contains a cycle through all its vertices, i.e., a connected spanning 2-regular subgraph.
- If  $H$  is a graph, then the **line graph of  $H$** , denoted by  $L(H)$ , is the graph on vertex set  $E(H)$  in which two distinct vertices in  $L(H)$  are adjacent if and only if their corresponding edges in  $H$  share an end vertex (with a straightforward extension in case of multiple edges).
- A graph is a **line graph** if it is isomorphic to  $L(H)$  for some graph  $H$ .
- Which graphs are line graphs and which are not?

# A forbidden subgraph characterization of line graphs

Theorem (Beineke, 1969)

A graph is a line graph if and only if it does not contain a copy of any of the following nine graphs as an **induced** subgraph.



## Forbidden induced subgraphs – claw-free graphs

- Let  $G$  be a graph and let  $S$  be a nonempty subset of  $V(G)$ . Then the **subgraph of  $G$  induced by  $S$** , denoted by  $G[S]$ , is the graph with vertex set  $S$ , and all edges of  $G$  with both end vertices in  $S$ .
- $H$  is an **induced subgraph** of  $G$  if it is induced in  $G$  by some subset of  $V(G)$ .
- $G$  is  **$H$ -free** if  $H$  is not an induced subgraph of  $G$ .
- In particular, a graph  $G$  is **claw-free** if  $G$  does not contain a copy of the **claw**  $K_{1,3}$  as an induced subgraph.
- Direct inspection of Beineke's result shows:

*every line graph is claw-free*

## The two earlier conjectures revisited

### Conjecture (MS)

*Every 4-connected claw-free graph is hamiltonian.*

### Conjecture (T)

*Every 4-connected line graph is hamiltonian.*

- Since line graphs are claw-free the first conjecture is **stronger** than the second one.
- Or are they **equivalent**? (A question Herbert Fleischner posed during the EIDMA workshop on Hamiltonicity of 2-tough graphs, Hotel Hölterhof, Enschede, November 19-24, 1995.)

## A useful tool: the closure

- To answer the question affirmatively, Zdeněk Ryjáček introduced a **closure concept** for claw-free graphs at the same workshop.
- It is based on adding edges to a graph  $G$  without destroying the (non)hamiltonicity (similar to the Bondy-Chvátal closure).
- The edges are added by looking at a vertex  $v$  and the subgraph of  $G$  induced by  $N(v)$ : the set of *neighbors* of  $v$ .
- If  $G[N(v)]$  is **connected** and not a complete graph, all edges are added to turn  $G[N(v)]$  into a **complete** graph.
- This procedure is repeated in the new graph, etc., until it is impossible to add any more edges.

## The two conjectures are equivalent

### Theorem (Ryjáček, 1997)

Let  $G$  be a **claw-free** graph. Then

- the closure  $\text{cl}(G)$  is uniquely determined,
- $\text{cl}(G)$  is hamiltonian if and only if  $G$  is hamiltonian,
- $\text{cl}(G)$  is the **line graph** of a triangle-free graph.

### Corollary (using a result of Zhan, 1991)

Every 7-connected claw-free graph is hamiltonian.

The conjectures are false for 3-connected graphs. The best positive result to date is the following by Tomáš Kaiser and Petr Vrána (2012).

### Theorem (Kaiser and Vrána, 2012)

Every **5-connected** claw-free graph with **minimum degree at least 6** is hamiltonian.

# From line graphs to their root graphs

- Whenever we consider a **line graph**  $G$ , we can identify a graph  $H$  such that  $G = L(H)$  (in polynomial time).
- If  $G$  is connected this  $H$  is unique, except for  $G = K_3$ : then  $H$  can be  $K_3$  or  $K_{1,3}$  (this is different for multigraphs).
- If we take  $K_{1,3}$  in this exceptional case, we can talk of a unique  $H$  as the **root graph** of the connected line graph  $G$  isomorphic to  $L(H)$ .
- What is the counterpart in  $H$  of a **Hamilton cycle** in  $G$ ?
- A *closed trail* (circuit) is a connected eulerian subgraph, i.e., a connected subgraph in which all degrees are even.
- A **dominating closed trail** (DCT or D-circuit) is a closed trail  $T$  such that every edge has at least one end vertex on  $T$ .

# Hamilton cycles and dominating closed trails

There is an intimate relationship between DCTs in  $H$  and Hamilton cycles in  $L(H)$ .

Theorem (Harary and Nash-Williams, 1965)

*Let  $H$  be a graph with at least three edges. Then  $L(H)$  is hamiltonian if and only if  $H$  contains a DCT.*

- What is the counterpart in  $H$  of 4-connectivity in  $L(H)$ ? Note that 4-edge-connectivity is not the right answer!
- A graph  $H$  is *essentially 4-edge-connected* if it contains no edge-cut  $R$  such that  $|R| < 4$  and at least two components of  $H - R$  contain an edge.
- $L(H)$  is 4-connected if and only if  $H$  is essentially 4-edge-connected.

## Another equivalent conjecture

The previous results and observations imply that the following conjecture is **equivalent** to the two we have seen before.

### Conjecture (DCT-conjecture)

Every **essentially** 4-edge-connected graph has a DCT.

- Note that 4-edge-connected graphs contain two edge-disjoint spanning trees.
- Hence 4-edge-connected graphs contain a **spanning** closed trail, in particular a DCT.
- So line graphs of 4-edge-connected graphs are hamiltonian (and *Hamilton-connected*).
- So the gap looks small .....

## From root graphs to cubic graphs

- If  $H$  is *cubic*, i.e., 3-regular, then a DCT becomes a *dominating cycle* (abbreviated DC).
- A cubic graph is essentially 4-edge-connected if and only if it is cyclically 4-edge-connected.
- $H$  is *cyclically 4-edge-connected* if  $H$  contains no edge-cut  $R$  such that  $|R| < 4$  and at least two components of  $H - R$  contain a cycle.

### Conjecture (Ash & Jackson, 1984)

Every cyclically 4-edge-connected **cubic** graph has a DC.

Fleischner and Jackson (1989) proved that this conjecture is **equivalent** to the others.

Main ingredient: Let  $H$  be an essentially 4-edge-connected graph of minimum degree  $\delta(G) \geq 3$  and let  $v \in V(H)$  be of degree  $d(v) \geq 4$ . Then some *cubic inflation* of  $H$  at  $v$  is essentially 4-edge-connected.

## Two weaker (?) conjectures from the same paper

### Conjecture (Jaeger, $\leq 1989$ )

*Every cyclically 4-edge-connected cubic graph  $G$  has a cycle  $C$  such that  $G - V(C)$  is acyclic.*

### Conjecture (Bondy, $\leq 1989$ )

*Every cyclically 4-edge-connected cubic graph  $G$  on  $n$  vertices has a cycle  $C$  of length at least  $c \cdot n$ , for some constant  $c$  with  $0 < c < 1$ .*

It is obvious that the conjecture of Ash-Jackson implies the conjecture of Jaeger, and one can show that the conjecture of Jaeger implies the conjecture of Bondy.

Are they equivalent?

# From cubic graphs to non-3-edge colorable cubic graphs

For *non-3-edge colorable* cubic graphs we have the following conjecture of Herbert Fleischner.

## Conjecture (F-Conjecture)

Every cyclically 4-edge-connected cubic graph that is **not 3-edge-colorable** has a DC.

Kochol (2000) proved that it is **equivalent** to the others.

One direction is obvious. For the other direction, he was assuming a counterexample to the previous conjecture and used it as a black box building block. In combination with an auxiliary gadget that is almost cubic and not 3-edge-colorable he constructed a counterexample to the F-Conjecture.

# From non-3-edge colorable cubic graphs to snarks

A *snark* is a cyclically 4-edge-connected cubic graph of **girth at least 5** that is not 3-edge-colorable.

## Conjecture (Snark-Conjecture)

*Every snark has a DC.*

The above conjecture is also **equivalent** to the others, as shown by B., Fijavž, Kaiser, Kužel, Ryjáček & Vrána (2008), using the constructive approach together with the concept of *contractible subgraphs*.

To date this is the **seemingly weakest** conjecture equivalent to the others.

Is there a link to the **Double Cycle Conjecture**?

Is there a link to **Nowhere Zero Flows**?

Let me turn to some **seemingly stronger** conjectures in the remainder of the talk.

## Seemingly stronger versions for cubic graphs

Fouquet & Thuillier (1990) established a seemingly stronger but equivalent version of the Ash-Jackson-Conjecture.

It is stronger in the sense that they require a DC that contains two given disjoint edges, as follows.

### Conjecture

*In a cyclically 4-edge-connected cubic graph **any two disjoint edges** are on a DC.*

The equivalence was extended by Fleischner & Kochol (2002) by requiring a DC through any two given edges.

### Conjecture

*In a cyclically 4-edge-connected cubic graph **any two edges** are on a DC.*

## Seemingly stronger versions for cubic graphs

There are several further equivalent versions involving **subgraphs** of cubic graphs.

### Conjecture (Kužel, 2008)

*Any subgraph  $H$  of an essentially 4-edge-connected cubic graph  $G$  with  $\delta(H) = 2$  and  $|V_2(H)| = 4$  is  $V_2(H)$ -dominated.*

### Conjecture (Kužel, Ryjáček & Vrána, 2012)

*Any subgraph  $H$  of an essentially 4-edge-connected cubic graph  $G$  with  $\delta(H) = 2$  and  $|V_2(H)| = 4$  is strongly  $V_2(H)$ -dominated.*

Here we assume that  $H$  has 4 vertices with degree 2, and we require that  $G + \{e, f\}$  has a DCT containing any two new (parallel) matching edges  $e$  and  $f$  added between pairs of these 4 vertices, and for the second statement also that  $G + e$  has a DCT containing any newly added edge  $e$  between a pair of these vertices.

## Seemingly stronger versions for cubic graphs

### Conjecture (B., Fijavž, Kaiser, Kužel, Ryjáček & Vrána, 2008)

*Every cyclically 4-edge-connected cubic graph  $G$  contains a weakly  $A_G(F)$ -contractible subgraph  $F$  with  $\delta(F) = 2$ .*

For a graph  $F$  and a set  $A \subset V(F)$ ,  $F$  is called (weakly)  $A$ -contractible if for every (nonempty) even subset  $X \subset A$  and for every partition  $P$  of  $X$  into 2-element subsets, the graph  $F^P$  has a DCT containing  $A$  and  $E(P)$ .

Here  $E(P) = \{uv \mid u, v \text{ are in the same equivalence class of } P\}$  and  $F^P$  is the multigraph obtained from  $F$  by adding the edges of  $E(P)$ .

If a 2-connected cubic graph  $H$  has a weakly  $A_H(F)$ -contractible subgraph  $F$ , then  $H$  has a DC if and only if  $H/F$  has a DCT.

Here  $A_H(F)$  are the vertices of  $F$  that have a neighbor in  $V(H) \setminus V(F)$ .

## Back to line graphs – seemingly stronger versions

A graph is *Hamilton-connected* if there is a Hamilton path between any two vertices.

Kužel & Xiong (2004) established equivalence with the following conjecture.

### Conjecture

*Every 4-connected line graph of a multigraph is **Hamilton-connected**.*

Ryjáček & Vrána (2011) extended the equivalence to claw-free graphs.

### Conjecture

*Every 4-connected claw-free graph is Hamilton-connected.*

Here the 2-closure played a crucial role: adding edges in the neighborhood if it induces a 2-connected graph. This does not always result in a line graph (of a multigraph): squares of cycles have to be treated differently.

## A link to the P versus NP problem

At present the seemingly strongest version of the conjectures is by Kužel, Ryjáček & Vrána (2012).

A graph  $G$  is *1-Hamilton-connected* if for any vertex  $x$  of  $G$  there is a Hamilton path in  $G - x$  between any two vertices. They also define a slightly stronger property called *2-edge-Hamilton-connectivity*.

### Conjecture

Every 4-connected line graph of a multigraph is **1-Hamilton-connected** (2-edge-Hamilton-connected).

The first statement can be extended to claw-free graphs. This is based on a further extension of the multigraph closure technique: the line graph obtained is of a multigraph with at most two triangles or at most one double edge.

This suggests that Thomassen's Conjecture (and all equivalent versions) might fail.

## How close are we to refuting the conjectures?

If the above conjecture is true, it implies that a line graph is 1-Hamilton-connected (2-edge-Hamilton-connected) **if and only if** it is 4-connected.

The connectivity of a (line) graph can be determined in **polynomial time**.

It is an **NP-complete** problem to decide whether a **line graph** is hamiltonian (Bertossi, 1981).

It is not difficult to show that deciding whether a given **graph** is 1-Hamilton-connected is also NP-complete.

It seems **likely** that deciding whether a given graph is 1-Hamilton-connected remains NP-complete when restricted to line graphs. If we can show this, it implies that

*Thomassen's Conjecture cannot be true, unless  $P=NP$ .*

## How close are we to proving the conjectures?

The gap between the conjecture(s) and the positive results is narrowing. If we drop the connectivity condition of the 2-regular spanning subgraph, we move from a Hamilton cycle to a *2-factor*. Enomoto, Jackson, Katerinis & Saito (1985) proved that every *2-tough* graph contains a 2-factor. This implies:

### Theorem

*Every 4-connected claw-free graph has a **2-factor**.*

It does not seem easy to use this as a starting point to show that there is a 2-factor with only one component, although there are some results that give upper bounds on the number of components. These results are beyond the scope of this talk.

## Relaxing the 4-connectedness and adding something else

I do not want to discuss degree conditions in this talk, but here is a connectivity-only result.

If we add an 'essentially connectivity' condition there is this result due to Lai, Shao, Wu & Zhou (2006).

### Theorem

*Every 3-connected, **essentially 11-connected** claw-free (line) graph is hamiltonian.*

Kaiser and Vrána claim that 11 can be replaced by 9, while 5 would be best possible (by the line graph of the Petersen graph in which the edges of a perfect matching are subdivided exactly once).

Question: how far can we decrease the 9 by raising the 3 to 4 in the theorem?

## Restrictions on the root graph

Lai (1994) proved the following partial affirmative answer to Thomassen's Conjecture.

### Theorem

*Every 4-connected line graph of a **planar** graph is hamiltonian.*

Kriesell (2001) proved a similar result on line graphs of claw-free (multi)graphs with the stronger conclusion of Hamilton-connectedness. In fact, he proved the following more general result.

### Theorem

*Let  $G$  be a graph such that  $L(G)$  is 4-connected and every vertex of degree 3 in  $G$  is on an edge of multiplicity at least 2 or on a triangle of  $G$ . Then  $L(G)$  is Hamilton-connected.*

## Restrictions on the root graph

For **quasi claw-free** graphs, i.e., in which all vertices  $u, v$  at distance 2 have a common neighbor  $w$  with  $N(w) \subseteq N[u] \cup N[v]$ , we have:

Theorem (Lai, Shao & Zhan, 2004)

*Every 4-connected line graph of a **quasi claw-free** graph is Hamilton-connected.*

There are many results along these lines.

In most proofs the root graphs are considered and the aim is to find a (closed or open) trail (internally) dominating all edges. A **common approach** is the following.

First the degree 1 vertices are deleted, then the degree 2 vertices are suppressed, and now one tries to show that the reduced graph has a suitable spanning (closed) trail. This is very similar to using two edge-disjoint spanning trees, or collapsibility, or advanced closure concepts.

# Bon appétit, guten Appetit, dobrou chut', eet smakelijk!

My favourite desserts for the end of the week:

- we will have improved on the currently best result that 5-connected claw-free graphs with  $\delta \geq 6$  are hamiltonian, or
- we will have established the complexity of 1-Hamilton-connectedness for line graphs, or
- we will have shown that all conjectures are equivalent to the famous Double Cycle Cover Conjecture, or
- ... this one is up to you!

I hope you did not lose your appetite!