# Hamilton-connected {claw,net}-free graphs, II

Xia Liu<sup>1,2</sup>, Zdeněk Ryjáček<sup>3,4,6</sup>, Petr Vrána<sup>3,4,6</sup>, Liming Xiong<sup>1,2,7</sup>, Xiaojing Yang<sup>1,2,5,8</sup>

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#### Abstract

In the first one in this series of two papers, we have proved that every 3-connected  $\{K_{1,3}, N_{1,3,3}\}$ -free graph is Hamilton-connected. In this paper, we continue in this direction by proving that every 3-connected  $\{K_{1,3}, X\}$ -free graph, where  $X \in \{N_{1,1,5}, N_{2,2,3}\}$ , is Hamilton-connected (where  $N_{i,j,k}$  is the graph obtained by attaching endvertices of three paths of lengths i, j, k to a triangle). This together with a previous result of other authors completes the characterization of forbidden induced generalized nets implying Hamilton-connectedness of a 3-connected claw-free graph. We also discuss remaining open cases in a full characterization of connected graphs X such that every 3-connected  $\{K_{1,3}, X\}$ -free graph is Hamilton-connected.

Keywords: Hamilton-connected; closure; forbidden subgraph; claw-free; net-free

### 1 Introduction

In this paper, we generally follow the most common graph-theoretical notations and terminology as given e.g. in [3], and for problem-specific notations and terminology we refer to the first paper of this series, [12]. We recall here the special graphs in Fig. 1 that will be important for our results.

The following two results were proved in [13] and in [12], respectively.

**Theorem A.** Let G be a 3-connected  $\{K_{1,3}, X\}$ -free graph, where

- (*i*) **[13]**  $X = N_{1,2,4}$ , or
- (*ii*) **[12]**  $X = N_{1,3,3}$ .

Then G is Hamilton-connected.

<sup>&</sup>lt;sup>1</sup>School of Mathematics and Statistics, Beijing Key Laboratory on MCAACI, Beijing Institute of Technology, Beijing 100081, P.R. China

<sup>&</sup>lt;sup>2</sup>e-mail liuxia\_90@163.com, lmxiong@bit.edu.cn, yangxiaojing89@163.com

<sup>&</sup>lt;sup>3</sup>Department of Mathematics; European Centre of Excellence NTIS - New Technologies for the Information Society, University of West Bohemia, Universitní 8, 301 00 Pilsen, Czech Republic

<sup>&</sup>lt;sup>4</sup>e-mail {ryjacek,vranap}@kma.zcu.cz

<sup>&</sup>lt;sup>5</sup>Department of Mathematics, Northwest Normal University, Lanzhou, Gansu 730070, P.R. China

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Figure 1: The graphs  $N_{i,j,k}$ ,  $S_{i,j,k}$  and the Wagner graph W

Recall that these results are sharp, as can be seen by considering the family of graphs  $\mathcal{G} = \{L(H) | H \in \mathcal{W}\}$ , where  $\mathcal{W}$  is the family of graphs obtained by attaching at least one pendant edge to each of the vertices of the Wagner graph W (see Fig. 1). Then any  $G \in \mathcal{G}$  is 3-connected, non-Hamilton-connected and  $N_{i,j,k}$ -free for  $i + j + k \geq 8$ . Hence the possible values that might imply a 3-connected  $\{K_{1,3}, N_{i,j,k}\}$ -free graph to be Hamilton-connected are those with  $i + j + k \leq 7$ . For i + j + k = 7, there are four possibilities, namely,  $N_{1,1,5}$ ,  $N_{1,2,4}$ ,  $N_{1,3,3}$  and  $N_{2,2,3}$ . While the second and third possibilities show sharpness of Theorem A, the first and last ones show that the next result, which is the main result of this paper, is also sharp, and completes the characterization.

**Theorem 1.** Let  $X \in \{N_{1,1,5}, N_{2,2,3}\}$ , and let G be a 3-connected  $\{K_{1,3}, X\}$ -free graph. Then G is Hamilton-connected.

The proof of Theorem 1, which is a careful case analysis, is postponed to Section 3. In Section 2, we collect necessary known results and facts that allow to significantly reduce the number of cases to be considered. Finally, in Section 4, we discuss sharpness and remaining open cases.

## 2 Preliminaries

In this section, we summarize some known facts that will be needed in our proof of Theorem 1. All these fact and results are contained already in the first paper of this series, [12], and we include them here for the sake of completeness.

### 2.1 Line graphs of multigraphs and their preimages

The line graph of a multigraph H is the graph G = L(H) with V(G) = E(H), in which two vertices are adjacent if and only if the corresponding edges of H share at least one vertex. It is well-known that in line graphs of multigraphs, for a given line graph G, a multigraph H such that G = L(H) is not uniquely determined. As shown in [16], this drawback can be overcome by an additional requirement that simplicial vertices correspond to pendant edges.

**Theorem B** [16]. Let G be a connected line graph of a multigraph. Then there is, up to an isomorphism, a uniquely determined multigraph H such that G = L(H) and a vertex  $e \in V(G)$  is simplicial in G if and only if the corresponding edge  $e \in E(H)$  is a pendant edge in H.

The multigraph H with the properties given in Theorem B will be called the *preimage* of a line graph G and denoted  $H = L^{-1}(G)$ . We will also use the notation a = L(e) and  $e = L^{-1}(a)$  for an edge  $e \in E(H)$  and the corresponding vertex  $a \in V(G)$ .

An edge-cut  $R \subset E(H)$  of a multigraph H is essential if H - R has at least two nontrivial components, and H is essentially k-edge-connected if every essential edge-cut of H is of size at least k. It is a well-known fact that a line graph G is k-connected if and only if  $L^{-1}(G)$ is essentially k-edge-connected. It is also a well-known fact that if X is a line graph, then a line graph G is X-free if and only if  $L^{-1}(G)$  does not contain as a subgraph (not necessarily induced) a graph F such that L(F) = X (but not necessarily  $F = L^{-1}(X)$ ). However, it is straightforward to verify that for the graph  $N_{i,j,k}$  there is exactly one graph F such that  $L(F) = N_{i,j,k}$ , namely, the graph  $L^{-1}(N_{i,j,k}) = S_{i+1,j+1,k+1}$  (see Fig. 1). Thus, we can conclude that a line graph G is  $N_{i,j,k}$ -free if and only if  $L^{-1}(G)$  does not contain as a (not necessarily induced) subgraph the graph  $L^{-1}(N_{i,j,k}) = S_{i+1,j+1,k+1}$ . Recall that when listing vertices of an  $S_{i,j,k}$  in a graph, we will write the list such that  $i \leq j \leq k$ , and we will use the notation  $S_{i,j,k}(v; a_1a_2 \ldots a_i; b_1b_2 \ldots b_j; c_1c_2 \ldots c_k)$  (in the labeling of vertices as in Fig. 1).

Recall that a closed trail T is a *dominating closed trail* (abbreviated DCT) if T dominates all edges of G, and an (e, f)-trail is an internally dominating (e, f)-trail (abbreviated (e, f)-IDT) if Int(T) dominates all edges of G. The following result shows the relation between a hamiltonian cycle (hamiltonian  $(a_1, a_2)$ -path) in G = L(H) and a DCT (an  $(e_1, e_2)$ -IDT) in H.

**Theorem C.** Let H be a multigraph with  $|E(H)| \ge 3$  and let G = L(H).

- (i) [8] The graph G is hamiltonian if and only if H has a DCT.
- (ii) [11] For every  $e_i \in E(H)$  and  $a_i = L(e_i)$ , i = 1, 2, G has a hamiltonian  $(a_1, a_2)$ -path if and only if H has an  $(e_1, e_2)$ -IDT.

#### 2.2 SM-closure

For  $x \in V(G)$ , the local completion of G at x is the graph  $G_x^* = (V(G), E(G) \cup \{y_1y_2 | y_1, y_2 \in N_G(x)\})$  (i.e.,  $G_x^*$  is obtained from G by adding all the missing edges with both vertices in  $N_G(x)$ ). Obviously, if G is claw-free, then so is  $G_x^*$ . Note that in the special case when G is a line graph and  $H = L^{-1}(G)$ , we have  $G_x^* = L(H|_e)$ , where  $e = L^{-1}(x)$ , and  $H|_e$  is obtained from H by contraction of e into a vertex and replacing the created loop(s) by pendant edge(s) (for more details on the contraction operation see Subsection 2.5). Finally, a vertex  $x \in V(G)$  is eligible if  $N_G(x)$  induces in G a connected noncomplete graph.

In [10], the concept of an SM-closure  $G^M$  of a claw-free graph G was defined by the following construction.

- (i) If G is Hamilton-connected, we set  $G^M = cl(G)$ .
- (*ii*) If G is not Hamilton-connected, we recursively perform the local completion operation at such eligible vertices for which the resulting graph is still not Hamilton-connected, as long as this is possible. We obtain a sequence of graphs  $G_1, \ldots, G_k$  such that
  - $G_1 = G$ ,
  - $G_{i+1} = (G_i)_{x_i}^*$  for some  $x_i \in V_{EL}(G_i), i = 1, \dots, k-1$ ,
  - $G_k$  has no hamiltonian (a, b)-path for some  $a, b \in V(G_k)$ ,
  - for any  $x \in V_{EL}(G_k)$ ,  $(G_k)_x^*$  is Hamilton-connected,

and we set  $G^M = G_k$ .

A resulting  $G^M$  is called a *strong M-closure* (or briefly an *SM-closure*) of the graph G, and a graph G equal to its SM-closure is said to be *SM-closed*. Note that for a given graph G, its SM-closure is not uniquely determined.

As shown in [16] and [10], if G is SM-closed, then G = L(H), where H does not contain any of the multigraphs shown in Fig. 2.



Figure 2: The diamond  $T_1$ , the multitriangle  $T_2$  and the triple edge  $T_3$ 

The following theorem summarizes basic properties of the SM-closure operation.

**Theorem D** [10]. Let G be a claw-free graph and let  $G^M$  be one of its SM-closures. Then  $G^M$  has the following properties:

- (i)  $V(G) = V(G^{\tilde{M}})$  and  $E(G) \subset E(G^{M})$ ,
- (ii)  $G^{M}$  is obtained from G by a sequence of local completions at eligible vertices,
- (iii) G is Hamilton-connected if and only if  $G^M$  is Hamilton-connected,
- (iv) if G is Hamilton-connected, then  $\tilde{G}^M = cl(G)$ ,
- (v) if G is not Hamilton-connected, then either
  - (a)  $V_{EL}(G^M) = \emptyset$  and  $G^M = \operatorname{cl}(G)$ , or
  - ( $\beta$ )  $V_{EL}(G^M) \neq \emptyset$  and  $(G^M)^*_{T}$  is Hamilton-connected for any  $x \in V_{EL}(G^M)$ ,
- (vi)  $G^M = L(H)$ , where H contains either
  - $(\alpha)$  at most 2 triangles and no multiedge, or
  - $(\beta)$  no triangle, at most one double edge and no other multiedge,
- (vii) if  $G^{M}$  contains no hamiltonian (a, b)-path for some  $a, b \in V(G^{M})$  and

  - (a) X is a triangle in H, then  $E(X) \cap \{L_{G^M}^{-1}(a), L_{G^M}^{-1}(b)\} \neq \emptyset$ , ( $\beta$ ) X is a multiedge in H, then  $E(X) = \{L_{G^M}^{-1}(a), L_{G^M}^{-1}(b)\}$ .

We will also need the following lemma on SM-closed graphs proved in [17].

Let G be an SM-closed graph, let  $H = L^{-1}(G)$  and let F be the Lemma E [17]. graph with  $V(F) = \{v_1, v_2, v_3, v_4, v_5, z\}$  and  $E(F) = \{v_1v_2, v_2v_3, v_3v_4, v_4v_5, v_5v_1, v_3v_5, zv_1, zv_2\}$ (see Fig. 3). Then H does not contain a subgraph  $\overline{H}$  isomorphic to the graph F such that  $N_H(\{v_1, v_2, v_3, v_5\}) \subset V(\bar{H}).$ 



Figure 3: The graph F

### 2.3 The core of the preimage of an SM-closed graph

The definition of the core is slightly problematic for multigraphs, therefore we restrict our observations to the case that we need, i.e., to preimages of 3-connected SM-closed graphs. The difficulties then do not occur since such a multigraph cannot have pendant multiedges by Theorem B, and cannot have pendant multitriangles (since there are no multitriangles at all).

Thus, let G be a 3-connected SM-closed graph and let  $H = L^{-1}(G)$ . The core of H is the multigraph co(H) obtained from H by removing all pendant edges and suppressing all vertices of degree 2.

Shao [20] proved the following properties of the core of a multigraph.

**Theorem F** [20]. Let H be an essentially 3-edge-connected multigraph. Then

- (i) co(H) is uniquely determined,
- (ii) co(H) is 3-edge-connected,
- (*iii*) V(co(H)) dominates all edges of H,
- (iv) if co(H) has a spanning closed trail, then H has a DCT.

#### 2.4 UM-closure

As shown in [12], the concept of SM-closure can be further strengthened by omitting the eligibility assumption in the local completion operation. Specifically, for a given claw-free graph G, we construct a graph  $G^U$  by the following construction.

- (i) If G is Hamilton-connected, we set  $G^U = K_{|V(G)|}$ .
- (*ii*) If G is not Hamilton-connected, we recursively perform the local completion operation at such vertices for which the resulting graph is still not Hamilton-connected, as long as this is possible. We obtain a sequence of graphs  $G_1, \ldots, G_k$  such that
  - $G_1 = G$ ,
  - $G_{i+1} = (G_i)_{x_i}^*$  for some  $x_i \in V(G_i), i = 1, \dots, k-1$ ,
  - $G_k$  has no hamiltonian (a, b)-path for some  $a, b \in V(G_k)$ ,
  - for any  $x \in V(G_k)$ ,  $(G_k)_x^*$  is Hamilton-connected,
  - and we set  $G^U = G_k$ .

A graph  $G^U$  obtained by the above construction is called an *ultimate M-closure* (or briefly a UM-closure) of the graph G, and a graph G equal to its UM-closure is said to be UM-closed.

Obviously, by the definition, if G is UM-closed, then G is also SM-closed, implying that G is a line graph and  $H = L^{-1}(G)$  has special structure (contains no diamond etc. – see Fig. 2 and Theorem D (vi), (vii)). The next theorem shows that for UM-closed graphs, not only H, but also co(H) has these strong structural properties.

**Theorem G** [12]. Let G be a claw-free graph and let  $G^U$  be one of its UM-closures. Then  $G^U$  has the following properties:

- (i)  $V(G) = V(G^U)$  and  $E(G) \subset E(G^U)$ ,
- (ii)  $G^U$  is obtained from G by a sequence of local completions at vertices,
- (*iii*) G is Hamilton-connected if and only if  $G^U$  is Hamilton-connected,
- (iv) if G is Hamilton-connected, then  $G^U = K_{|V(G)|}$ ,
- (v) if G is not Hamilton-connected, then  $(G^U)_x^*$  is Hamilton-connected for any  $x \in V(G^U)$ ,

- (vi)  $G^U = L(H)$ , where co(H) contains no diamond, no mutitriangle and no triple edge, and either
  - $(\alpha)$  at most 2 triangles and no multiedge, or
  - $(\beta)$  no triangle, at most one double edge and no other multiedge, and if co(H)contains a double edge, then this double edge is also in H,

(vii) if  $G^U$  contains no hamiltonian (a, b)-path for some  $a, b \in V(G^U)$  and

- (a) X is a triangle in co(H), then  $E(X) \cap \{L_{G^U}^{-1}(a), L_{G^U}^{-1}(b)\} \neq \emptyset$ , ( $\beta$ ) X is a multiedge in co(H), then  $E(X) = \{L_{G^U}^{-1}(a), L_{G^U}^{-1}(b)\}$ .

The following result was first established in [5], and later on reconsidered in [14] in a more general setting.

**Theorem H** [14]. Let G be a  $\{K_{1,3}, N_{i,j,k}\}$ -free graph,  $i, j, k \ge 1$ , and let  $x \in V(G)$ . Then the graph  $G_x^*$  is  $\{K_{1,3}, N_{i,j,k}\}$ -free.

Specifically, Theorem H implies that a UM-closure of a  $\{K_{1,3}, N_{i,j,k}\}$ -free graph is also  $\{K_{1,3}, N_{i,j,k}\}$ -free.

#### A-contractible multigraphs 2.5

For a multigraph H and  $F \subset H$ ,  $H|_F$  denotes the multigraph obtained from H by identifying the vertices of F as a (new) vertex  $v_F$ , and by replacing the created loops by pendant edges. Specifically, if  $E(F) = \{e\}$ , we simply write  $G|_e$ . If H is a multigraph,  $X \subset V(H)$ , and A is a partition of X into subsets, then  $E(\mathcal{A})$  denotes the set of all edges  $a_1a_2$  (not necessarily in H) such that  $a_1, a_2$  are in the same element of  $\mathcal{A}$ . Further  $H^{\mathcal{A}}$  denotes the multigraph with vertex set  $V(H^{\mathcal{A}}) = V(H)$  and edge set  $E(H^{\mathcal{A}}) = E(H) \cup E(\mathcal{A})$  (where E(H) and  $E(\mathcal{A})$  are considered to be disjoint, i.e., if  $e_1 = a_1 a_2 \in E(H)$  and  $e_2 = a_1 a_2 \in E(\mathcal{A})$ , then  $e_1, e_2$  are parallel edges in  $H^{\mathcal{A}}$ ).

Let F be a multigraph and let  $A \subset V(F)$ . We say that F is A-contractible, if for every even subset  $X \subset A$  and for every partition  $\mathcal{A}$  of X into two-element subsets, the graph  $F^{\mathcal{A}}$ has a DCT containing all vertices of A and all edges of  $E(\mathcal{A})$ . Note that this definition allows X to be empty, in which case  $F^{\mathcal{A}} = F$ . Also, if F is A-contractible, then F is A'-contractible for any  $A' \subset A$  (since every subset X of A' is a subset of A). As shown in [15], if F is an  $A_H(F)$ -contractible submultigraph of a multigraph H, then H has a DCT if and only if  $H|_F$ has a DCT (note that the concept was defined in [15] for graphs, but it is easy to observe that it remains true also for multigraphs).

Several examples of A-contractible graphs are shown in Fig. 4 (where the vertices in the set A are double-circled). Note that detailed proofs of A-contractibility are for  $F_2$  and  $F_3$  given in [15].

In our proof, we will need the following lemma from [12].

Let H be a multigraph, F an  $A_H(F)$ -contractible submultigraph of H, and Lemma I [12]. let  $e_1, e_2 \in E(H) \setminus E(F)$ . Then H has an  $(e_1, e_2)$ -IDT if and only if  $H|_F$  has an  $(e_1, e_2)$ -IDT.



Figure 4: Examples of A-contractible graphs

### 2.6 A special version of the "Nine-point-theorem"

The well-known "Nine point theorem" by Holton et al. [9] states that a 3-connected cubic graph contains a cycle passing through any 9 prescribed vertices. For our proof, we will need its special version, based on a stronger version by Bau and Holton [1], and developed in [12].

**Theorem J** [12]. Let  $X \in \{N_{1,1,5}, N_{1,3,3}, N_{2,2,3}\}$ , and let G be a 3-connected UM-closed  $\{K_{1,3}, X\}$ -free graph such that  $\operatorname{co}(H)$ , where  $H = L^{-1}(G)$ , is 2-connected. Let  $e_1, e_2 \in E(H)$  be such that there is no  $(e_1, e_2)$ -IDT in H. Then for every set  $A \subset V(\operatorname{co}(H))$ , |A| = 8, there is an  $(e_1, e_2)$ -trail T in H such that  $A \subset \operatorname{Int}(T)$ .

The following lemma will be crucial in our proof. Recall that W denotes the Wagner graph (see Fig. 1).

**Lemma K** [12]. Let G be a 3-connected non-Hamilton-connected UM-closed claw-free graph. Then G has an induced subgraph  $\tilde{G}$  (possibly  $\tilde{G} = G$ ) such that  $\tilde{G}$  is 3-connected, non-Hamilton-connected and UM-closed, and, moreover,  $\tilde{H}_0 = \operatorname{co}(L^{-1}(\tilde{G}))$  is 2-connected, and either  $c(\tilde{H}_0) \geq 9$  and  $|V(\tilde{H}_0)| \geq 10$ , or  $\tilde{H}_0 \simeq W$ .

### 3 Proof of Theorem 1

Let G be a 3-connected  $\{K_{1,3}, N_{1,1,5}\}$ -free or  $\{K_{1,3}, N_{2,2,3}\}$ -free graph and suppose, to the contrary, that G is not Hamilton-connected. By Theorem G and Theorem H, we can suppose that G is UM-closed. Let thus  $H = L^{-1}(G)$ , and set  $H_0 = co(H)$ . By Theorem F(ii),  $H_0$  is 3-edge-connected. By Lemma K, we can suppose that  $H_0$  is 2-connected and  $c(H_0) \ge 9$  and  $|V(H_0)| \ge 10$ , unless  $H_0 \simeq W$ . Then, by Theorems J and C(ii), we have the following claim.

<u>Claim 1.</u> Let  $A \subset V(H_0)$  be such that |A| = 8. Then A does not dominate all edges of H.

<u>Proof.</u> Since G is not Hamilton-connected, by Theorem C(ii), there are edges  $e_1, e_2 \in E(H)$  such that there is no  $(e_1, e_2)$ -IDT in H. Then, by Theorem J, there is an  $(e_1, e_2)$ -trail T in H such that  $A \subset Int(T)$ . But if A dominates all the edges in H, then T would be an  $(e_1, e_2)$ -IDT in H.

Now, if  $H_0 \simeq W$ , then  $|V(H_0)| = 8$  and  $V(H_0)$  dominates all edges of H, contradicting Claim 1. Thus, we have  $c(H_0) \ge 9$  and  $|V(H_0)| \ge 10$ . We consider separate cases for possible values of  $c(H_0)$  and  $|V(H_0)|$ .

Throughout the proof, in each of the cases,  $C = x_1 x_2 \dots x_{c(H_0)}$  always denotes a longest cycle in  $H_0$ ,  $R = V(H) \setminus V(C)$ ,  $N = \{y \in V(H_0) | N_R(y) = \emptyset\}$ ,  $R_0 = R \cap V(H_0)$ , and if  $R_0 \neq \emptyset$ , we set  $R_0 = \{y_1, \dots, y_{|R_0|}\}$  and we choose the notation such that  $y_1 x_1 \in E(H_0)$ . An edge  $x_i x_j$  with  $x_i, x_j \in V(C)$ ,  $1 \leq i, j \leq |V(C)|$ , will be called a *chord* of C, and we say that  $x_i x_j$  is a *k*-chord if the shorter one of the two subpaths of C determined by  $x_i$  and  $x_j$  has kinterior vertices. Similarly, for a vertex  $y \in R_0$  with  $N_C(y) \neq \emptyset$ , the subpaths of C determined by the neighbors of y will be called *segments* of C, and a segment with k interior vertices will be referred to as a *k*-segment.

<u>Claim 2.</u> If  $E(\langle R \rangle_H) = \emptyset$ , then  $H_0$  has no double edge.

<u>Proof.</u> Suppose that  $H_0$  has a double edge  $\{e, f\}$ . Then, since V(C) dominates all edges of H, it follows that H has an (e, f)-IDT. But then, by Theorem  $G(vii)(\beta)$ , G has an (a, b)-path for every pair  $a, b \in V(G)$ , a contradiction.

In the proof, we will often list vertices of a subgraph  $S_{i,j,k}$ . There are two general comments to all these situations.

- When some edge  $e = x_i x_j$  of the  $S_{i,j,k}$  is in  $E(H_0)$ , it can always happen that e is subdivided in H, i.e., formally,  $e \notin E(H)$ . However, it is immediate to see that if this happens, then the corresponding subgraph of H, which instead of  $e = x_i x_j$  contains a path  $x_i z x_j$  with  $z \in V_2(H)$ , also contains  $S_{i,j,k}$  as a subgraph.
- When a vertex  $x_i \in V(C)$  has a (potential) neighbor  $z \in R$  and the vertex z occurs as the last vertex of a branch of the  $S_{i,j,k}$ , then such a vertex z can be an endvertex of a pendant edge attached to  $x_i$ , or can be  $z \in V_2(H)$  and z subdivides some of the edges incident to  $x_i$ . It should be noted that in the second case, the vertices  $x_i$  and z can occur in reverse order in the list (i.e.,  $x_i$  being the last vertex of the branch).

Throughout the proof, we always implicitly understand that there are also these possibilities.

<u>Case 1</u>: G is  $\{K_{1,3}, N_{1,1,5}\}$ -free.

Then H does not contain as a subgraph the graph  $S_{2,2,6}$ .

**Subcase 1.1:**  $c(H_0) = 9$  and  $|V(H_0)| \ge 10$ .

First observe that  $E(\langle R \rangle_H) = \emptyset$ , since if e.g.  $y_1 z \in E(H)$  for some  $z \in R$ , then H contains the subgraph  $S_{2,2,6}(x_1; y_1 z; x_2 x_3; x_9 x_8 x_7 x_6 x_5 x_4)$ , a contradiction. Consequently, no vertex in  $R_0$  is connected to C by a double edge by Claim 2. Thus,  $y_1$  has at least three distinct neighbors on C since  $d_{H_0}(y_1) \geq 3$ .

Suppose that  $y_1x_3 \in E(H_0)$ . Then  $x_2 \in N$ , for otherwise, if  $x_2z \in E(H)$  for some  $z \in R$ , either there is  $S_{2,2,6}(x_1; x_2z; y_1x_3; x_9x_8x_7x_6x_5x_4)$  in H if  $z \neq y_1$ , or the cycle  $C' = x_1y_1x_2\ldots x_9x_1$  is longer than C if  $z = y_1$ , a contradiction. Hence  $\{x_2, y_1\} \subset N$  and, by the same argument for  $y_2, \ldots, y_{|R_0|}$  and since  $E(\langle R \rangle_H) = \emptyset$ , we have  $\{x_2, y_1, \ldots, y_{|R_0|}\} \subset N$ . This implies that the set  $A = \{x_1, x_3, x_4, x_5, x_6, x_7, x_8, x_9\}$  with |A| = 8 dominates all edges in H, contradicting Claim 1. Thus,  $y_1x_3 \notin E(H_0)$ .

Since  $d_C(y_1) \geq 3$  and |V(C)| = 9, the only possibility is, up to a symmetry, that  $N_C(y_1) = \{x_1, x_4, x_7\}$ . Then again  $x_2 \in N$ , for otherwise, for a neighbor z of  $x_2$  in R, we have  $S_{2,2,6}(x_1; x_2z; x_9x_8; y_1x_7x_6x_5x_4x_3)$  in H if  $z \neq y_1$ , or a longer cycle if  $z = y_1$ , a contradiction. By symmetry, we have  $V(C) \setminus \{x_1, x_4, x_7\} \subset N$ , implying that, if  $|V(H_0)| > 10$ ,

all  $y_i$ ,  $i = 2, ..., |R_0|$ , are adjacent to  $x_1$ ,  $x_4$  and  $x_7$ , and to no other vertices of C. Then the set  $A = V(C) \setminus \{x_9\}$  with |A| = 8 dominates all edges in H, contradicting Claim 1.

#### **Subcase 1.2:** $c(H_0) = |V(H_0)| = 10.$

By Theorem G(vi),  $H_0$  has at most two triangles. Hence, since  $\delta(H_0) \ge 3$ , it follows that C has a k-chord for some  $k \ge 2$ . We consider possible cases.

#### <u>Subcase 1.2.1</u>: $x_1x_4 \in E(H_0)$ .

We show that  $\{x_2, x_5, x_7\} \subset N$ . Thus, let  $z \in R$  and consider the following possibilities.

Edge	$S_{2,2,6}$ in $H$
$x_2 z \in E(H)$	$S_{2,2,6}(x_1; x_2 z; x_4 x_3; x_{10} x_9 x_8 x_7 x_6 x_5)$
$x_5 z \in E(H)$	$S_{2,2,6}(x_4; x_3x_2; x_5z; x_1x_{10}x_9x_8x_7x_6)$
$x_7 z \in E(H)$	$S_{2,2,6}(x_4; x_3x_2; x_5x_6; x_1x_{10}x_9x_8x_7z)$

Hence we have  $\{x_2, x_5, x_7\} \subset N$  and, by symmetry, also  $\{x_3, x_8, x_{10}\} \subset N$ . Now, the set  $A = V(C) \setminus \{x_3, x_5\}$  with |A| = 8 dominates all edges of H, contradicting Claim 1, unless  $x_3x_5 \in E(H_0)$ . Hence  $x_3x_5 \in E(H_0)$ . Symmetrically,  $x_2x_{10} \in E(H_0)$ . Similarly, considering the set  $A = V(C) \setminus \{x_5, x_7\}$  with |A| = 8, we have  $x_5x_7 \in E(H_0)$ . But then  $x_3x_5, x_2x_{10}$  and  $x_5x_7$  are three 1-chords in C, determining three triangles in  $H_0$ , which contradicts Theorem G(vi). Thus, C has no 2-chord.

#### <u>Subcase 1.2.2</u>: $x_1x_5 \in E(H_0)$ .

Then  $x_6 \in N$ , for otherwise we have a subgraph  $S_{2,2,6}(x_1; x_2x_3; x_5x_4; x_{10}x_9x_8x_7x_6z)$  in H for  $z \in R$  with  $x_6z \in E(H)$ . Hence  $x_6 \in N$  and, symmetrically,  $x_{10} \in N$ . If  $x_6x_{10} \notin E(H)$ , then  $A = V(C) \setminus \{x_6, x_{10}\}$  is a set of 8 vertices in  $H_0$  that dominates all edges of H, contradicting Claim 1. Thus,  $x_6x_{10} \in E(H_0)$ . By symmetry,  $\{x_1, x_5, x_6, x_{10}\} \subset N$ . Then, considering the set  $A = V(C) \setminus \{x_1, x_6\}$ , Claim 1 implies  $x_1x_6 \in E(H_0)$ . Symmetrically,  $x_5x_{10} \in E(H_0)$ . But then  $\langle \{x_1, x_5, x_6, x_{10}\} \rangle_H \simeq K_4$ , contradicting Theorem G(vi). Thus, C has no 3-chord.

By Subcases 1.2.1 and 1.2.2 and since |V(C)| = 10, C can have only 4-chords, plus at most two 1-chords. Thus, at most 4 vertices of C can be in a 1-chord, implying that C has at least three 4-chords. Since |V(C)| = 10, there always is a pair of 4-chords such that their endvertices are on C at distance 2. Thus, let, say,  $x_1x_6 \in E(H_0)$  and  $x_3x_8 \in E(H_0)$ . Then  $x_2 \in N$ , since otherwise, for a  $z \in R$  with  $x_2z \in E(H)$ , we have  $S_{2,2,6}(x_1; x_2z; x_{10}x_9; x_6x_7x_8x_3x_4x_5)$  in H. Symmetrically also  $x_7 \in N$ . Considering the set  $A = V(C) \setminus \{x_2, x_7\}$  with |A| = 8, by Claim 1 we have  $x_2x_7 \in E(H_0)$ . Now,  $x_1 \in N$ , since otherwise, for a  $z \in R$  with  $x_1z \in E(H)$  we have  $S_{2,2,6}(x_6; x_1z; x_7x_2; x_5x_4x_3x_8x_9x_{10})$  in H. Thus, we have  $\{x_1, x_2, x_7\} \subset N$ . Considering the set  $A = V(C) \setminus \{x_1, x_7\}$  with |A| = 8, we have  $x_1x_7 \in E(H_0)$  by Claim 1. However,  $x_1x_7$  is a 3-chord, and, by symmetry, we are back in Subcase 1.2.2.

#### **Subcase 1.3:** $c(H_0) = 10$ and $|V(H_0)| \ge 11$ .

Then  $E(\langle R \rangle_H) = \emptyset$ , for if there is a  $z \in R$  with e.g.  $zy_1 \in E(H)$ , then H contains  $S_{2,2,6}(x_1; y_1z; x_2x_3; x_{10}x_9x_8x_7x_6x_5)$ . Hence  $R \subset N$ , and no vertex in  $R_0$  is attached to C by a double edge by Claim 2. Since  $\delta(H_0) \geq 3$ , every vertex in  $R_0$  has at least three distinct neighbors on C. We consider the possible cases.

<u>Subcase 1.3.1</u>:  $y_1x_4 \in E(H_0)$ .

Then we have  $S_{2,2,6}(x_1; y_1x_4; x_2x_3; x_{10}x_9x_8x_7x_6x_5)$  in  $H_0$ , a contradiction.

<u>Subcase 1.3.2</u>:  $y_1x_5 \in E(H_0)$ .

Then  $x_6 \in N$  since otherwise H contains  $S_{2,2,6}(x_1; y_1x_5; x_2x_3; x_{10}x_9x_8x_7x_6z)$  for some  $z \in R$  with  $zx_6 \in E(H)$ . Symmetrically also  $x_{10} \in N$ . If  $x_6x_{10} \in E(H_0)$ , then the set  $A = V(C) \setminus \{x_6, x_{10}\}$  with |A| = 8 dominates all edges of H (recall that  $R \subset N$ ). Hence  $x_6x_{10} \in E(H_0)$  by Claim 1. But then H contains  $S_{2,2,6}(x_6; x_7x_8; x_{10}x_9; x_5y_1x_1x_2x_3x_4)$ , a contradiction.

By Subcases 1.3.1 and 1.3.2, by symmetry and since C is longest, the neighbors of  $y_1$  on C can be at distance (along C) either 2, or at least 5. However, it is straightforward to verify that this is not possible since |V(C)| = 10.

<u>Subcase 1.4</u>:  $c(H_0) = |V(H_0)| = 11.$ 

Since  $\delta(H_0) \geq 3$  and  $R_0 = \emptyset$ , every vertex of C is in a chord. We consider the possible cases.

<u>Subcase 1.4.1</u>:  $x_1x_4 \in E(H_0)$ .

Then we have  $S_{2,2,6}(x_1; x_2x_3; x_4x_5; x_{11}x_{10}x_9x_8x_7x_6)$  in  $H_0$ , a contradiction. Thus, by symmetry, there is no 2-chord.

<u>Subcase 1.4.2</u>:  $x_1x_5 \in E(H_0)$ .

Then we have  $S_{2,2,6}(x_1; x_2x_3; x_5x_4; x_{11}x_{10}x_9x_8x_7x_6)$  in  $H_0$ , a contradiction. Thus, by symmetry, there is no 3-chord.

We observe the following fact.

<u>Claim 3.</u> If  $e = x_i x_{i+5}$  is a 4-chord of C, then  $\{x_{i-1}, x_{i+6}\} \subset N$  (indices modulo 11).

<u>Proof.</u> Choose the notation such that i = 1. If  $x_7$  is adjacent to a  $z \in R$ , then we have  $S_{2,2,6}(x_1; x_2x_3; x_6x_5; x_{11}x_{10}x_9x_8x_7z)$  in H, hence  $x_7 \in N$ . Symmetrically,  $x_{11} \in N$ .

We show that C has no 1-chord. Thus, let, to the contrary,  $x_1x_3 \in E(H_0)$ . By Theorem G(vi),  $H_0$  contains no diamond, and hence  $x_4, x_{11} \notin N_{H_0}(x_2)$ . Since  $\delta(H_0) \geq 3$ and C has neither 2-chords nor 3-chords, we may assume, by symmetry, that  $x_2x_7 \in E(H_0)$ . By Claim 3, we have  $\{x_1, x_8\} \subset N$ . Moreover,  $x_2 \in N$ , for otherwise we have  $S_{2,2,6}(x_1; x_2z; x_3x_4; x_{11}x_{10}x_9x_8x_7x_6)$  in H for some neighbor z of  $x_2$  in R, and also  $x_{10} \in N$ since for a  $z \in R$  with  $zx_{10} \in E(H)$  we have  $S_{2,2,6}(x_7; x_8x_9; x_6x_5; x_2x_3x_1x_{11}x_{10}z)$  in H. Thus, we have  $\{x_1, x_2, x_8, x_{10}\} \subset N$ .

Considering the set  $A = V(C) \setminus \{x_2, x_8, x_{10}\}$  with |A| = 8 and since  $x_2x_{10} \notin E(H_0)$  by Subcase 2.4.1, either  $x_2x_8 \in E(H_0)$  or  $x_8x_{10} \in E(H_0)$  by Claim 1. Similarly, considering the set  $A = V(C) \setminus \{x_1, x_8, x_{10}\}$  with |A| = 8 and since  $x_1x_8 \notin E(H_0)$  by Subcase 2.4.2, either  $x_1x_{10} \in E(H_0)$  or  $x_8x_{10} \in E(H_0)$ . Thus, there is at least one of the edges  $x_2x_8$ ,  $x_8x_{10}$ , and also at least one of the edges  $x_1x_{10}, x_8x_{10}$ . Since each of these edges creates a triangle, and there already is one triangle in  $H_0$  (created by the 1-chord  $x_1x_3$ ), the only possibility is that  $x_8x_{10} \in E(H_0)$ . Now, since  $d_{H_0}(x_{11}) \geq 3$ , by the previous subcases and since  $H_0$  does not contain a diamond, we have  $x_{11}x_5 \in E(H_0)$  or  $x_{11}x_6 \in E(H_0)$  (recall that we also have  $\{x_1, x_2, x_8, x_{10}\} \subset N$ ).

Let first  $x_{11}x_5 \in E(H_0)$ . Then  $x_6 \in N$  by Claim 3, and considering the set  $A = V(C) \setminus \{x_2, x_6, x_{10}\}$  with |A| = 8, Claim 1 implies that  $H_0$  contains at least one of the edges  $x_2x_6, x_6x_{10}, x_2x_{10}$ . However, each of these edges contradicts Subcase 1.4.1 or 1.4.2. Thus, we have  $x_{11}x_6 \in E(H_0)$ . Then similarly  $x_5 \in N$  by Claim 3, and we have an analogous contradiction by Claim 1 considering the set  $A = V(C) \setminus \{x_1, x_5, x_8\}$ . Thus, C has no 1-chord.

By the previous considerations, C has no k-chord for k = 1, 2, 3. Since  $\delta(H_0) \geq 3$ , every vertex of C is in a 4-chord. By symmetry, let  $x_1x_6 \in E(H_0)$ . Then  $x_9$  is in a 4-chord and, again by symmetry, we can suppose that  $x_3x_9 \in E(H_0)$ . Then  $H_0$  contains  $S_{2,2,6}(x_9; x_8x_7; x_{10}x_{11}; x_3x_2x_1x_6x_5x_4)$ , a contradiction.

**Subcase 1.5:**  $c(H_0) \ge 11$  and  $|V(H_0)| > c(H_0)$ .

First observe that, as in Subcase 1.3,  $E(\langle R \rangle_H) = \emptyset$  (otherwise there is an  $S_{2,2,6}$  in H), implying  $R \subset N$ , and that no vertex in  $R_0$  is attached to C by a double edge by Claim 2. Since  $\delta(H_0) \geq 3$ , every vertex in  $R_0$  has at least three distinct neighbors on C, and we consider possible cases. Set |V(C)| = t.

<u>Subcase 1.5.1</u>:  $y_1x_4 \in E(H_0)$ .

Then we have  $S_{2,2,6}(x_1; x_2x_3; y_1x_4; x_tx_{t-1}x_{t-2}x_{t-3}x_{t-4}x_{t-5})$ . Thus, there is no 2-segment on C.

<u>Subcase 1.5.2</u>:  $y_1x_5 \in E(H_0)$ .

Then we similarly have  $S_{2,2,6}(x_1; x_2x_3; y_1x_5; x_tx_{t-1}x_{t-2}x_{t-3}x_{t-4}x_{t-5})$ . Thus, there is no 3-segment on C.

By Subcases 1.5.1 and 1.5.2,  $y_1$  determines on C only 1-segments and k-segments for  $k \ge 4$ , and there are at least three segments determined by  $y_1$ . We distinguish possible cases, where, in each of the cases, we always consider  $k_1$ -segments and  $k_2$ -segments such that  $k_1 \le k_2$  and  $k_1 + k_2$  is smallest possible.

<u>Subcase 1.5.3</u>:  $k_1 = k_2 = 1$ .

By symmetry, we can choose the notation such that  $\{x_1, x_3, x_5\} \subset N_C(y_1)$ , and then we have  $S_{2,2,6}(y_1; x_1x_2; x_3x_4; x_5x_6x_7x_8x_9x_{10})$  in H.

<u>Subcase 1.5.4</u>:  $k_1 = 1, k_2 \ge 4$ .

Then, by symmetry,  $\{x_1, x_3, x_{k_2+4}\} \subset N_C(y_1)$  and, by the choice of  $k_1$  and  $k_2$ ,  $t \geq 2k_2 + 4 \geq 12$ . Then we have  $S_{2,2,6}(x_3; x_2x_1; x_4x_5; y_1x_{k_2+4}x_{k_2+5}x_{k_2+6}x_{k_2+7}x_{k_2+8})$  in H.

<u>Subcase 1.5.5</u>:  $4 \le k_1 \le k_2$ .

Then  $\{x_1, x_{k_1+2}, x_{k_1+k_2+3}\} \subset N_C(y_1)$  and  $t \ge k_1 + 2k_2 + 3 \ge 15$ . Then we have  $S_{2,2,6}(x_1; x_2x_3; x_tx_{t-1}; y_1x_{k_1+2}x_{k_1+3}x_{k_1+4}x_{k_1+5}x_{k_1+6})$ .

#### **Subcase 1.6:** $c(H_0) = |V(H_0)| \ge 12.$

Since  $\delta(H_0) \ge 3$  and  $R_0 = \emptyset$ , every vertex of C is in a chord, and we consider the possible cases. Set |V(C)| = t.

<u>Subcase 1.6.1</u>:  $x_1x_4 \in E(H_0)$ .

Then we have  $S_{2,2,6}(x_1; x_2x_3; x_4x_5; x_tx_{t-1}x_{t-2}x_{t-3}x_{t-4}x_{t-5})$  in  $H_0$ , a contradiction. By symmetry, there is no 2-chord.

**Subcase 1.6.2:**  $x_1x_i \in E(H_0), i \in \{5, 6\}.$ 

Then we have  $S_{2,2,6}(x_1; x_2x_3; x_ix_{i-1}; x_tx_{t-1}x_{t-2}x_{t-3}x_{t-4}x_{t-5})$  in  $H_0$ , a contradiction. Thus, by symmetry, there is no 3-chord and no 4-chord.

<u>Subcase 1.6.3</u>:  $x_1x_7 \in E(H_0)$ .

If  $t \geq 13$ , we similarly have  $S_{2,2,6}(x_1; x_2x_3; x_7x_6; x_tx_{t-1}x_{t-2}x_{t-3}x_{t-4}x_{t-5})$ . Thus, suppose that t = 12. By the previous subcases, there are only 1-chords and 5-chords. Since  $H_0$ has at most two triangles, there are at most two 1-chords, and since t = 12 and every vertex of C is in a chord, there are three vertices that are consecutive on C and each of them is in a 5-chord. Choose the notation such that  $x_1x_7 \in E(H_0), x_2x_8 \in E(H_0)$ and  $x_3x_9 \in E(H_0)$ . Then we have  $S_{2,2,6}(x_1; x_2x_8; x_{12}x_{11}; x_7x_6x_5x_4x_3x_9)$  in  $H_0$ .

<u>Subcase 1.6.4</u>:  $x_1x_i \in E(H_0), i \ge 8$ .

Then, by the definition of a chord (the shorter one of the two paths),  $t \ge 2i - 2 \ge 14$ , and we have  $S_{2,2,6}(x_1; x_i x_{i+1}; x_t x_{t-1}; x_2 x_3 x_4 x_5 x_6 x_7)$  in  $H_0$ .

<u>**Case 2:**</u> G is  $\{K_{1,3}, N_{2,2,3}\}$ -free.

Then H does not contain as a subgraph the graph  $S_{3,3,4}$ . In this case, we make an additional choice of the cycle C; namely, we choose C such that

(i) C is a longest cycle in  $H_0$ , and

(ii) subject to (i), C dominates the maximum number of edges of H.

**Subcase 2.1:**  $c(H_0) = 9$  and  $|V(H_0)| \ge 10$ .

Subcase 2.1.1:  $E(\langle R \rangle_H) \neq \emptyset$ .

Let  $e \in E(\langle R \rangle_H)$ . By Theorem F(*iii*), we can choose the notation such that  $e = y_1 z_1$ , where  $y_1 \in R_0$  and  $z_1 \in R$ , and, if possible, we choose e and  $y_1$  such that  $|N_C(y_1)| \ge 3$ .

<u>Subcase 2.1.1.1:</u>  $|N_C(y_1)| \ge 3.$ 

We first show that  $x_5, x_6 \notin N_{H_0}(y_1)$ . Let, say, e.g.  $y_1x_5 \in E(H_0)$ . Clearly  $x_2, x_4, x_6, x_9 \notin N_C(y_1)$  since C is longest. By the assumption of the subcase and by symmetry,  $y_1x_3 \in E(H_0)$  or  $y_1x_7 \in E(H_0)$ . If  $y_1x_3 \in E(H_0)$ , then  $N_R(x_4) = \emptyset$ , since otherwise, for a  $z \in N_R(x_4)$  we have  $S_{3,3,4}(x_1; y_1x_5x_6; x_9x_8x_7; x_2x_3x_4z)$  in H, but then the cycle  $C' = x_1x_2x_3y_1x_5\ldots x_9x_1$  dominates in H more edges than C, contradicting choice (*ii*) of C. Similarly, if  $y_1x_7 \in E(H_0)$ , then  $N_R(x_6) = \emptyset$  (otherwise we have  $S_{3,3,4}(x_1; x_2x_3x_4; x_9x_8x_7; y_1x_5x_6z)$  in H for a  $z \in N_R(x_6)$ ), and then the cycle  $C' = x_1x_2x_3x_4x_5y_1x_7x_8x_9x_1$  dominates in H more edges than C, contradicting choice (*ii*) of C. Thus, by symmetry,  $x_5, x_6 \notin N_{H_0}(y_1)$ .

Since  $|N_C(y_1)| \ge 3$ ,  $y_1$  determines on C three segments, say,  $k_i$ -segments, i = 1, 2, 3, and since  $|V(C) \setminus N_C(y_1)| \le 6$ , the possible distributions of interior vertices of the segments are  $(k_1, k_2, k_3) = (1, 1, 4)$ ,  $(k_1, k_2, k_3) = (1, 2, 3)$  or  $(k_1, k_2, k_3) = (2, 2, 2)$  (where we admit that the 4-segment in the case (1, 1, 4) or the 3-segment in the case (1, 2, 3) can be further subdivided by another neighbor of  $y_1$  if  $|N_C(y_1)| > 3$ ).

However, in the first case we have  $y_1x_5 \in E(H_0)$  and in the second case we have  $y_1x_6 \in E(H_0)$  (or a symmetric situation), a contradiction. Thus, we have  $(k_1, k_2, k_3) = (2, 2, 2)$ , implying  $N_C(y_1) = \{x_1, x_4, x_7\}$ .

We observe that  $N_R(x_i) = \emptyset$  for i = 2, 3, 5, 6, 8, 9 since if, say,  $y_2 z \in E(H)$  for some  $z \in R$ , we have  $S_{3,3,4}(x_4; x_3x_2z; x_5x_6x_7; y_1x_1x_9x_8)$  in H; other cases are symmetric. Since  $H_0$  is 3-edge-connected, each of the vertices  $x_2, x_3, x_5, x_6, x_8, x_9$  is either in a double edge or in a chord. By Theorem G(vi) and by symmetry, we can choose the notation such the segment  $x_1x_2x_3x_4$  contains neither a double edge nor a 1-chord. Considering possible chords containing  $x_2$ , we have the following possibilities.

Chord	Contradiction
$x_2 x_5$	$C' = x_1 y_1 x_4 x_3 x_2 x_5 x_6 x_7 x_8 x_9 x_1$ is longer than C
$x_{2}x_{6}$	$S_{3,3,4}(x_2; x_3x_4x_5; x_1y_1z_1; x_6x_7x_8x_9)$ in $H$
$x_2 x_7$	$S_{3,3,4}(x_2; x_7y_1z_1; x_1x_9x_8; x_3x_4x_5x_6)$ in $H$
$x_2 x_8$	$C' = x_1 y_1 x_7 x_6 x_5 x_4 x_3 x_2 x_8 x_9 x_1$ is longer than C
$x_{2}x_{9}$	$S_{3,3,4}(x_2; x_1y_1z_1; x_3x_4x_5; x_9x_8x_7x_6)$ in $H$

In each of the cases, we have reached a contradiction.

#### **Subcase 2.1.1.2:** $|N_C(y_1)| = 2$ .

By the connectivity, one of the two connections of  $y_1$  to C is a double edge and hence, by Theorem  $G(vi)(\beta)$ , there is no other double edge in  $H_0$ . Hence  $y_1$  is the only vertex in  $H_0$  that has a neighbor in R. This also implies that  $H_0$  is trianglefree. Obviously,  $x_2, x_9 \notin N_C(y_1)$  since C is longest. Thus, by symmetry, we have  $y_1x_3 \in E(H_0), y_1x_4 \in E(H_0)$  or  $y_1x_5 \in E(H_0)$ . In each of these cases, each of the vertices in  $V(C) \setminus N_C(y_1)$  (specifically,  $x_2$  in the first two cases and  $x_3$  in the third case) must be connected to some other vertex of C by a chord or by a path P of length two with an interior vertex in  $R_0$  since  $H_0$  is 3-edge-connected.

Below we list (up to a symmetry) all possible cases. Here, if we give an  $S_{3,3,4}$  for some chord, say,  $x_2x_j$ , it is always implicitly understood that  $x_2x_j$  can also be a path of length two, in which case the listed graph contains an  $S_{3,3,4}$  as a proper subgraph. Also note that although there is no 1-chord since  $H_0$  is triangle-free,  $x_2$  can be still connected to  $x_4$  or to  $x_9$  by a path of length two with interior vertex, say,  $y_2 \in R_0$ .

Case	Chord or path	Contradiction: $S_{3,3,4}$ in $H$
$y_1 x_3 \in E(H_0)$	$P = x_2 y_2 x_4$	$S_{3,3,4}(x_2; y_2 x_4 x_5; x_3 y_1 z_1; x_1 x_9 x_8 x_7)$
	$x_2x_5 \in E(H_0)$	$  S_{3,3,4}(x_5; x_2x_1x_9; x_6x_7x_8; x_4x_3y_1z_1)  $
	$x_2 x_6 \in E(H_0)$	$S_{3,3,4}(x_2; x_3x_4x_5; x_1y_1z_1; x_6x_7x_8x_9)$
$y_1 x_4 \in E(H_0)$	$P = x_2 y_2 x_4$	$S_{3,3,4}(x_4; y_2 x_2 x_3; y_1 x_1 x_9; x_5 x_6 x_7 x_8)$
	$x_2x_5 \in E(H_0)$	$  S_{3,3,4}(x_2; x_5x_6x_7; x_1x_9x_8; x_3x_4y_1z_1)  $
	$x_2 x_6 \in E(H_0)$	$  S_{3,3,4}(x_2; x_3x_4x_5; x_1y_1z_1; x_6x_7x_8x_9)  $
	$x_2x_7 \in E(H_0)$	$  S_{3,3,4}(x_2; x_7 x_8 x_9; x_1 y_1 z_1; x_3 x_4 x_5 x_6)  $
	$x_2 x_8 \in E(H_0)$	$  S_{3,3,4}(x_8; x_2x_3x_4; x_7x_6x_5; x_9x_1y_1z_1)  $
	$P = x_2 y_2 x_9$	$S_{3,3,4}(x_2; x_3x_4x_5; x_1y_1z_1; y_2x_9x_8x_7)$
$y_1 x_5 \in E(H_0)$	$P = x_3 y_2 x_5$	$S_{3,3,4}(x_5; y_2 x_3 x_4; y_1 x_1 x_2; x_6 x_7 x_8 x_9)$
	$x_3x_6 \in E(H_0)$	$  S_{3,3,4}(x_3; x_6x_7x_8; x_2x_1x_9; x_4x_5y_1z_1)  $
	$x_3x_7 \in E(H_0)$	$S_{3,3,4}(x_3; x_4x_5x_6; x_7x_8x_9; x_2x_1y_1z_1)$

In each of the possible cases, we have reached a contradiction.

#### Subcase 2.1.2: $E(\langle R \rangle_H) = \emptyset$ .

By Claim 2,  $H_0$  has no double edge, implying that every vertex in  $R_0$  has three distinct neighbors on C. If, say,  $y_1x_5 \in E(H_0)$ , then  $x_4$  has no neighbor in R (otherwise, for a  $z \in N_R(x_4)$ , we have  $S_{3,3,4}(x_1; y_1x_5x_6; x_9x_8x_7; x_2x_3x_4z)$  in H), and then, considering the set  $A = V(C) \setminus \{x_4\}$  with |A| = 8, we have a contradiction by Claim 1. Hence, by symmetry,  $y_1$  adjacent to neither  $x_5$  nor  $x_6$ . This implies that, similarly as in Subcase 2.1.1.1, for  $y_1$ ,  $(k_1, k_2, k_3) = (2, 2, 2)$  is the only possible distribution of segments on C. Thus, we have  $N_C(y_1) = \{x_1, x_4, x_7\}$ . If there is a  $z \in N_R(x_3)$ , we have  $S_{3,3,4}(x_1; x_2x_3z; y_1x_4x_5; x_9x_8x_7x_6)$  in H, hence  $N_R(x_3) = \emptyset$ . But then again, considering the set  $A = V(C) \setminus \{x_3\}$  with |A| = 8, we have a contradiction by Claim 1.

**Subcase 2.2:**  $c(H_0) = |V(H_0)| = 10.$ 

Subcase 2.2.1: C has a 3-chord.

Let  $x_1x_5 \in E(H_0)$  be a 3-chord of C, and let  $z \in R$  be adjacent (in H) to a vertex on C. We consider the following cases.

Case	Contrdiction: $S_{3,3,4}$ in $H$
$zx_2 \in E(H)$	$S_{3,3,4}(x_5; x_6x_7x_8; x_1x_{10}x_9; x_4x_3x_2z)$
$zx_6 \in E(H)$	$S_{3,3,4}(x_1; x_2x_3x_4; x_5x_6z; x_{10}x_9x_8x_7)$
$zx_7 \in E(H)$	$S_{3,3,4}(x_5; x_4x_3x_2; x_6x_7z; x_1x_{10}x_9x_8)$

Thus, we have  $N_R(\{x_2, x_6, x_7\}) = \emptyset$ , and, by symmetry, also  $N_R(\{x_4, x_9, x_{10}\}) = \emptyset$ . Now, considering the set  $A_1 = V(C) \setminus \{x_2, x_4\}$  with  $|A_1| = 8$ , we have  $x_2x_4 \in E(H_0)$  by Claim 1. Similarly, considering the set  $A_2 = V(C) \setminus \{x_7, x_9\}$  with  $|A_2| = 8$  and the set  $A_3 = V(C) \setminus \{x_2, x_{10}\}$  with  $|A_3| = 8$ , we have  $x_7x_9 \in E(H_0)$  and  $x_2x_{10} \in E(H_0)$  by Claim 1. But then  $T_1 = x_2x_3x_4x_2$ ,  $T_2 = x_7x_8x_9x_7$  and  $T_3 = x_2x_1x_{10}x_2$  are three triangles in  $H_0$ , contradicting Theorem G(v). Thus, C has no 3-chord.

#### Subcase 2.2.2: C has a 4-chord.

Let  $x_1x_6 \in E(H_0)$  be a 4-chord of C. Then  $N_R(x_3) = \emptyset$  (otherwise, for a  $z \in N_R(x_3)$ , we have  $S_{3,3,4}(x_1; x_2x_3z; x_6x_5x_4; x_{10}x_9x_8x_7)$  in H), and, symmetrically,  $N_R(x_9) = \emptyset$ . Moreover,  $x_3x_9 \notin E(H_0)$  by Subcase 2.2.1. But then, considering the set  $A = V(C) \setminus \{x_3, x_9\}$  with |A| = 8, we have a contradiction by Claim 1. Thus, C has no 4-chord.

By Subcases 2.2.1 and 2.2.2., every vertex of C is in a 1-chord or in a 2-chord. Since there are at most two 1-chords by Theorem G(v), we can choose the notation such that  $x_1$  is in a 2-chord, i.e.,  $x_1x_4 \in E(H_0)$ .

<u>Claim 4.</u> If  $x_i x_{i+3}$  is a 2-chord of C, then  $N_R(\{x_{i+1}, x_{i+2}\}) = \emptyset$  (indices modulo 10).

<u>Proof.</u> By symmetry, set i = 1. If  $z \in N_R(x_3)$ , then  $S_{3,3,4}(x_1; x_2x_3z; x_4x_5x_6; x_{10}x_9x_8x_7)$  is a subgraph of H, a contradiction. Hence  $N_R(x_3) = \emptyset$ . Symmetrically,  $N_R(x_2) = \emptyset$ .  $\Box$ 

Thus, by Claim 4, we have  $N_R(\{x_2, x_3\}) = \emptyset$ . Suppose that  $x_2$  is in a 2-chord. Then either  $x_2x_5 \in E(H_0)$ , or  $x_2x_9 \in E(H_0)$ .

Let first  $x_2x_5 \in E(H_0)$ . Then we have  $N_R(x_4) = \emptyset$  by Claim 4, and, considering the set  $A = V(C) \setminus \{x_2, x_4\}$  with |A| = 8, we have  $x_2x_4 \in E(H_0)$  by Claim 1. But then  $\langle \{x_1, x_2, x_4, x_5\} \rangle_{H_0}$  is a diamond in  $H_0$ , a contradiction. Hence  $x_2x_5 \notin E(H_0)$ .

Secondly, let  $x_2x_9 \in E(H_0)$ . Then, again by Claim 4, we have  $N_R(\{x_1, x_{10}\}) = \emptyset$ . Considering the set  $A = V(C) \setminus \{x_2, x_{10}\}$  with |A| = 8, we have  $x_2x_{10} \in E(H_0)$  by Claim 1, but then  $\langle \{x_1, x_2, x_9, x_{10}\} \rangle_{H_0}$  is a diamond or a  $K_4$  in  $H_0$ , a contradiction again. Hence  $x_2$  is in a 1-chord. Symmetrically,  $x_3$  is also in a 1-chord, and since  $H_0$  is diamondfree, we have  $x_2x_{10} \in E(H_0)$  and  $x_3x_5 \in E(H_0)$ . Then  $N_R(x_4) = \emptyset$ , since otherwise we have  $S_{3,3,4}(x_2; x_1x_4z; x_3x_5x_6; x_{10}x_9x_8x_7)$  in H for a  $z \in N_R(x_4)$ . Recall that also  $N_R(x_2) = \emptyset$  by Claim 4. But then, considering the set  $A = V(C) \setminus \{x_2, x_4\}$  with |A| = 8, we have  $x_2x_4 \in E(H_0)$ , implying that  $\langle \{x_1, x_2, x_3, x_4\} \rangle_{H_0}$  is a diamond or a  $K_4$  in  $H_0$ , a contradiction.

**Subcase 2.3:**  $c(H_0) \ge 10$  and  $|V(H_0)| > c(H_0)$ .

Set |V(C)| = t. If  $y_1$  determines on C a 3-segment, i.e.,  $y_1x_5 \in E(H_0)$ , we have  $S_{3,3,4}(x_1; x_2x_3x_4; y_1x_5x_6; x_tx_{t-1}x_{t-2}x_{t-3})$  in  $H_0$ , if  $y_1$  determines on C a 4-segment, i.e.,  $y_1x_6 \in E(H_0)$ , we have  $S_{3,3,4}(x_1; y_1x_6x_7; x_tx_{t-1}x_{t-2}; x_2x_3x_4x_5)$  in  $H_0$ , and if  $y_1$  determines on C a k-segment with  $k \geq 5$ , i.e.,  $y_1x_{k+2} \in E(H_0)$ , then, by the definition of a segment (the shorter one of the two subpaths of C), necessarily  $t \geq 2k + 2$ , and then we have  $S_{3,3,4}(x_1; y_1x_{k+2}x_{k+3}; x_tx_{t-1}x_{t-2}; x_2x_3x_4x_5)$  in  $H_0$ . Thus, by symmetry, every vertex in  $R_0$  determines on C only 1-segments and 2-segments.

Now, if a vertex in  $R_0$ , say,  $y_1$ , has three distinct neighbors on C, then necessarily  $y_1$  is adjacent to some of  $x_5$ ,  $x_6$  or  $x_7$ , and we are in some of the previous cases. Thus,  $y_1$  has two neighbors on C, and one of the two connections of  $y_1$  to C is a double edge. This implies that  $R_0 = \{y_1\}, |V(H_0)| = t + 1$ , and  $H_0$  is triangle-free. Choose the notation such that  $N_C(y_1) = \{x_1, x_3\}$  or  $N_C(y_1) = \{x_1, x_4\}$ . Since  $\delta(H_0) \ge 3$ , the vertex  $x_2$  has, besides  $x_1$  and  $x_3$ , another neighbor in  $H_0$ . Clearly  $N_{R_0}(x_2) = \emptyset$  since  $R_0 = \{y_1\}$  and Cis longest, and also  $x_2x_4, x_2x_t \notin E(H_0)$  since  $H_0$  is triangle-free.

Let first  $y_1x_3 \in E(H_0)$ . Up to a symmetry, we have the following possibilities (note that if  $x_2x_k \in E(H_0)$  for  $k \ge 7$ , then  $t \ge 2k - 4$  by symmetry).

Case	Contradiction: $S_{3,3,4}$ in $H_0$
$x_2x_5 \in E(H_0)$	$S_{3,3,4}(x_1; y_1x_3x_4; x_2x_5x_6; x_tx_{t-1}x_{t-2}x_{t-3})$
$x_2 x_6 \in E(H_0)$	$S_{3,3,4}(x_1; y_1x_3x_4; x_2x_6x_5; x_tx_{t-1}x_{t-2}x_{t-3})$
$x_2 x_k \in E(H_0),  k \ge 7$	$S_{3,3,4}(x_1; y_1x_3x_4; x_tx_{t-1}x_{t-2}; x_2x_kx_{k-1}x_{k-2})$

Thus,  $y_1 x_4 \in E(H_0)$ . Then we have the following possibilities.

Case	Contradiction: $S_{3,3,4}$ in $H_0$
$x_2 x_5 \in E(H_0)$	$S_{3,3,4}(x_1; y_1x_4x_3; x_2x_5x_6; x_tx_{t-1}x_{t-2}x_{t-3})$
$x_2 x_6 \in E(H_0)$	$S_{3,3,4}(x_1; y_1x_4x_3; x_2x_6x_5; x_tx_{t-1}x_{t-2}x_{t-3})$
$x_2 x_k \in E(H_0), \ 7 \le k \le t - 3$	$S_{3,3,4}(x_1; y_1x_4x_3; x_tx_{t-1}x_{t-2}; x_2x_kx_{k-1}x_{k-2})$
$x_2 x_{t-2} \in E(H_0)$	$  S_{3,3,4}(x_2; x_3x_4y_1; x_1x_tx_{t-1}; x_{t-2}x_{t-3}x_{t-4}x_{t-5})  $
$x_2 x_{t-1} \in E(H_0)$	$S_{3,3,4}(x_4; y_1x_1x_t; x_3x_2x_{t-1}; x_5x_6x_7x_8)$

**Subcase 2.4:**  $c(H_0) = |V(H_0)| \ge 11$ .

Set |V(C)| = t. We recall that  $H_0$  has no double edge by Claim 2. Since  $\delta(H_0) \ge 3$ , every vertex of C is in at least one chord.

If C has a 3-chord, say,  $x_1x_5 \in E(H_0)$ , we have  $S_{3,3,4}(x_1; x_2x_3x_4; x_5x_6x_7; x_tx_{t-1}x_{t-2}x_{t-3})$ in  $H_0$ , and, similarly, if C has a k-chord for  $k \ge 4$ , then  $t \ge \max\{2k+2, 11\}$  by the definition of a chord, and we have  $S_{3,3,4}(x_1; x_{k+2}x_{k+3}x_{k+4}; x_tx_{t-1}x_{t-2}; x_2x_3x_4x_5)$  in  $H_0$ . Thus, the only possible chords in C are 1-chords and 2-chords.

We show that there is no 1-chord. Let, to the contrary,  $x_1x_3 \in E(H_0)$  be a 1-chord of C. Since  $x_2$  must be in a chord and a 1-chord at  $x_2$  would create a diamond, by symmetry, we have  $x_2x_5 \in E(H_0)$ . We consider possible chords at  $x_4$ . Since both  $x_1x_4$  and  $x_2x_4$ create a diamond, necessarily  $x_4x_6 \in E(H_0)$  or  $x_4x_7 \in E(H_0)$ . However, if  $x_4x_7 \in E(H_0)$ , we have  $S_{3,3,4}(x_1; x_2x_5x_6; x_3x_4x_7; x_tx_{t-1}x_{t-2}x_{t-3})$  in  $H_0$ , hence  $x_4x_6 \in E(H_0)$ . Now, if  $x_5$ has another neighbor on C, then the only possibility is  $x_5x_8 \in E(H_0)$  (all other chords at  $x_5$  create a diamond), but then we have  $S_{3,3,4}(x_1; x_3x_4x_6; x_tx_{t-1}x_{t-2}; x_2x_5x_8x_7)$  in  $H_0$ . Thus,  $x_5x_8 \notin E(H_0)$ , and, symmetrically,  $x_2x_{10} \notin E(H_0)$ .

We summarize that the subgraph  $F = \langle \{x_1, x_2, x_3, x_4, x_5, x_6\} \rangle_{H_0}$  is isomorphic to the graph F in Fig. 3, and the vertices  $x_2, x_3, x_4, x_5$  have no other neighbors on C outside F. Moreover, if, say,  $x_5$  has a neighbor  $z \in R$ , we have  $S_{3,3,4}(x_1; x_2x_5z; x_3x_4x_6; x_tx_{t-1}x_{t-2}x_{t-3})$  in H; thus, by symmetry, none of the vertices  $x_2, x_3, x_4, x_5$  has a neighbor in R. We conclude that F is a subgraph of H, and  $N_H(\{x_2, x_3, x_4, x_5\}) \subset V(F)$ , which contradicts Lemma E. Thus, C has no 1-chord, implying that the only possible chords are 2-chords.

Claim 5. If  $x_i x_{i+3} \in E(H_0)$  is a 2-chord of C for some  $i, 1 \leq i \leq t$ , then  $d_{H_0}(x_{i+1}) = d_{H_0}(x_{i+2}) = 3$  (indices modulo t).

<u>Proof.</u> Clearly  $3 \leq d_{H_0}(x) \leq 4$  for  $x \in V(C)$ . Let, say,  $x_3x_6 \in E(H_0)$  be a 2chord of C, and suppose that  $d_{H_0}(x_4) = 4$ . Then  $x_1x_4, x_4x_7 \in E(H_0)$ , and we have  $S_{3,3,4}(x_1; x_4x_7x_8; x_tx_{t-1}x_{t-2}; x_2x_3x_6x_5)$  in  $H_0$ . Hence  $d_{H_0}(x_4) = 3$ , and, symmetrically,  $d_{H_0}(x_5) = 3$ . <u>Claim 6.</u> Let  $x_i x_{i+3}$ ,  $x_{i+1} x_{i+4}$ ,  $x_{i+2} x_{i+5}$  be 2-chords of C for some  $i, 1 \le i \le t$  (indices modulo t), and let  $X = \langle \{x_i, x_{i+1}, x_{i+2}, x_{i+3}, x_{i+4}, x_{i+5}\} \rangle_{H_0}$ . Then

- (i) X is isomorphic to the graph  $F_4$  in Fig. 4,
- (*ii*)  $N_H(\{x_{i+1}, x_{i+2}, x_{i+3}, x_{i+4}\}) \subset V(X),$
- (*iii*)  $A_{H_0}(X) = \{x_i, x_{i+5}\}.$

<u>Proof.</u> (i) Part (i) is obvious since any further edge in X would be a k-chord of C for  $k \neq 2$ .

(*ii*) Let, say,  $x_{i+2}y \in E(H)$  for some  $y \notin V(X)$ . If  $y \in V(C)$ , then  $d_{H_0}(x_{i+2}) = 4$ , contradicting Claim 5, and if  $y \in R$ , we have  $S_{3,3,4}(x_i; x_{i+1}x_{i+2}y; x_{i+3}x_{i+4}x_{i+5}; x_{i-1}x_{i-2}x_{i-3}x_{i-4})$  in H. Thus,  $N_H(x_{i+2}) \subset V(X)$ , and, by symmetry,  $N_H(\{x_{i+1}, x_{i+2}, x_{i+3}, x_{i+4}\}) \subset V(X)$ . (*iii*) By (*ii*),  $x_i$  and  $x_{i+5}$  are the only vertices of X having a neighbor outside X.

We show that  $H_0$  has no vertices of degree 4. Let, to the contrary, say,  $d_{H_0}(x_6) = 4$ . Then  $x_6$  is in two 2-chords, i.e.,  $x_3x_6, x_6x_9 \in E(H_0)$ . By Claim 5,  $d_{H_0}(x_4) = d_{H_0}(x_5) = 3$ , implying  $|N_{H_0}(x_4) \cap \{x_1, x_7\}| = |N_{H_0}(x_5) \cap \{x_2, x_8\}| = 1$ . If  $x_4x_7 \in E(H_0)$ , then  $H_0$  contains  $S_{3,3,4}(x_6; x_3x_2x_1; x_9x_{10}x_{11}; x_5x_4x_7x_8)$ ; hence  $x_1x_4 \in E(H_0)$ . By symmetry,  $x_5x_8 \notin E(H_0)$ , implying  $x_2x_5 \in E(H_0)$ . Symmetrically,  $x_7x_{10}, x_8x_{11} \in E(H_0)$ . Set  $X_1 = \langle \{x_1, x_2, x_3, x_4, x_5, x_6\} \rangle_{H_0}$  and  $X_2 = \langle \{x_6, x_7, x_8, x_9, x_{10}, x_{11}\} \rangle_{H_0}$ . By Claim 6, we have  $X_1 \simeq X_2 \simeq F_4$  (where  $F_4$  is the graph in Fig. 4),  $A_{H_0}(X_1) = \{x_1, x_6\}, A_{H_0}(X_2) = \{x_6, x_{11}\}$ , and  $X_i$  is  $A_{H_0}(X_i)$ -contractible, i = 1, 2.

Let first t = 11. The set  $\{x_3, x_5, x_9\}$  is independent in  $H_0$  and, by Claim 6(ii), we have  $N_R(\{x_3, x_5, x_9\}) = \emptyset$ . Considering the set  $A = V(C) \setminus \{x_3, x_5, x_9\}$  with |A| = 8, we have a contradiction by Claim 1. Thus,  $t \ge 12$ .

Then necessarily  $d_{H_0}(x_1) = 4$  or  $d_{H_0}(x_{11}) = 4$ , for otherwise  $\{x_1x_t, x_{11}x_{12}\}$  is an edge-cut of  $H_0$ , a contradiction. By symmetry, let  $d_{H_0}(x_{11}) = 4$ . Then  $x_{11}$  is, besides the chord  $x_8x_{11}$ , in another 2-chord of C, implying  $t \ge 14$  and  $x_{11}x_{14} \in E(H_0)$ . However, since  $x_{12}$  and  $x_{13}$  must have a chord, by Claim 5, we have  $t \ge 16$  and  $x_{12}x_{15}, x_{13}x_{16} \in E(H_0)$ . Set  $X_3 = \langle \{x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16} \} \rangle_{H_0}$ . By Claim 6, we have  $X_3 \simeq F_4$ ,  $N_H(\{x_{12}, x_{13}, x_{14}, x_{15}\}) \subset V(X_3)$ ,  $A_{H_0}(X_3) = \{x_{11}, x_{16}\}$ , and  $X_3$  is  $A_{H_0}(X_3)$ -contractible. But then  $X_1, X_2$  and  $X_3$  are three contractible subgraphs of H, contradicting Lemma I. Thus, there is no vertex of degree 4, i.e.,  $H_0$  is cubic, implying that t is even and  $t \ge 12$ .

Let now  $x_1x_4$  be a 2-chord of C. Since  $d_{H_0}(x_2) = d_{H_0}(x_3) = 3$ , up to a symmetry, either  $x_2x_5, x_3x_6 \in E(H_0)$ , or  $x_2x_{t-1}, x_3x_6 \in E(H_0)$  (or  $x_2x_5, x_3x_t \in E(H_0)$ , which is symmetric with the first possibility). If  $x_2x_5, x_3x_6 \in E(H_0)$ , then, by Claim 6, again  $X_1 = \langle \{x_1, x_2, x_3, x_4, x_5, x_6\} \rangle_{H_0} \simeq F_4$ ,  $N_H(\{x_2, x_3, x_4, x_5\}) \subset V(X_1)$  and  $A_{H_0}(X_1) = \{x_1, x_6\}$ , but then  $\{x_1x_t, x_6x_7\}$  is an edge-cut of  $H_0$ , a contradiction. Thus,  $x_2x_{t-1}, x_3x_6 \in E(H_0)$ . Now,  $x_2x_5 \notin E(H_0)$  by Claim 5, and since  $d_{H_0}(x_5) = 3, x_5x_8 \in E(H_0)$ . But then we have  $S_{3,3,4}(x_1; x_4x_5x_8; x_tx_{t-1}x_{t-2}; x_2x_3x_6x_7)$  in  $H_0$ , a contradiction.

### 4 Concluding remarks

1. Theorems A and 1 can be slightly extended as follows. For  $s \ge 0$ , a graph G is s-Hamiltonconnected if the graph G - M is Hamilton-connected for any set  $M \subset V(G)$  with  $|M| \le s$ . Obviously, an s-Hamilton-connected graph must be (s+3)-connected. Since an induced subgraph of a  $\{K_{1,3}, N_{i,j,k}\}$ -free graph is also  $\{K_{1,3}, N_{i,j,k}\}$ -free, we immediately have the following fact, showing that the obvious necessary condition is also sufficient in  $\{K_{1,3}, N_{i,j,k}\}$ -free graphs.

**Corollary 2.** Let s, i, j, k be integers such that  $s \ge 0$ ,  $i, j, k \ge 1$  and  $i + j + k \le 7$ , and let G be a  $\{K_{1,3}, N_{i,j,k}\}$ -free graph. Then G is s-Hamilton-connected if and only if G is (s+3)-connected.

2. We will now discuss sharpness of the known results and the remaining open cases in the characterization of all pairs of connected graphs X, Y that might imply a 3-connected  $\{X, Y\}$ -free graph to be Hamilton-connected. To avoid trivial cases, we restrict the observations to  $X, Y \not\simeq P_3$ . Starting from the negative side, we recall the following result that appeared in [4].

**Theorem L** [4]. If X, Y is a pair of connected graphs such that  $X, Y \not\simeq P_3$  and every 3-connected  $\{X, Y\}$ -free graph is Hamilton connected, then, up to a symmetry,  $X = K_{1,3}$  and Y satisfies each of the following conditions:

- (a)  $\Delta(Y) \leq 3$ ,
- (b) any longest induced path in Y has at most 9 vertices,
- (c) Y contains no cycles of length at least 4,
- (d) the distance between two distinct triangles in Y is either 1 or at least 3,
- (e) There are at most two triangles in Y,
- (f) Y is claw-free.

Moreover, item (d) was reduced in [7] to read

(d') the distance between two distinct triangles in Y is either 1 or 3.

The reduction in [7] consists in two steps: excluding even lengths, and showing that length 5 is not possible (knowing that length more than 5 is not possible by (b)). The graph used in the first step can be easily turned into an infinite family by attaching in the preimage arbitrarily many pendant edges to vertices of degree 3; however, the graph used in the second step has 20 vertices and any additional pendant edge or subdivision in the preimage makes it contain an induced  $\Gamma_5$ . Thus, distance 5 might be still possible for  $n = |V(G)| \ge 21$ . We therefore replace (d') with the following more precise statement:

(d'') the distance between two distinct triangles in Y is 1 or 3, or possibly also 5 for  $n = |V(G)| \ge 21$ .

This implies that, as noted in [2], the only possibilities for the graph Y are (see Fig. 5 for the graphs  $Z_i$ ,  $B_{i,j}$ ,  $N_{i,j,k}$  and  $\Gamma_i$ ):

- (i) the path  $P_i$  with  $4 \le i \le 9$ ,
- (*ii*) the graph  $Z_i$ ,
- (*iii*) the generalized bull  $B_{i,j}$ ,
- (iv) the generalized net  $N_{i,j,k}$ ,
- (v) the generalized hourglass  $\Gamma_1$  or  $\Gamma_3$ , or also  $\Gamma_5$  for  $n = |V(G)| \ge 21$ ,

- (vi) the generalized hourglass  $\Gamma_1$  or  $\Gamma_3$ , or also  $\Gamma_5$  for  $n = |V(G)| \ge 21$ , with paths possibly attached to either of the two triangles.
- In [2], this list was further reduced by excluding the possibility (vi).

Note that the reduction in [2], and also all observations in the proof of Theorem L in [4], are based on infinite families of graphs. Thus, except for the  $\Gamma_5$ , there is no hope to get a corresponding result for "sufficiently large" graphs.



Figure 5: The graphs  $\Gamma_i$ ,  $Z_i$ ,  $B_{i,j}$  and  $N_{i,j,k}$ 

To obtain upper bounds on the possible values of the parameters i, j, k, consider the family of graphs  $\mathcal{G} = \{L(H) | H \in \mathcal{W}\}$ , introduced in Section 1. The graphs in  $\mathcal{G}$  are 3-connected, non-Hamilton-connected,  $P_{10}$ -free,  $Z_i$ -free for  $i \geq 7$  and with a single exception of the smallest graph in  $\mathcal{G}$  (having 20 vertices) even  $Z_i$ -free for  $i \geq 8$ ,  $B_{i,j}$ -free for  $i + j \geq 8$ , and  $N_{i,j,k}$ -free for  $i + j + k \geq 8$ . Hence the possible graphs Y that might imply a 3-connected  $\{K_{1,3}, Y\}$ -free graph to be Hamilton-connected are  $\Gamma_1$ ,  $\Gamma_3$ ,  $\Gamma_5$  for  $n \geq 21$ ,  $P_i$  for  $i \leq 9$ ,  $Z_i$  for  $i \leq 6$  and  $Z_7$ for  $n \geq 21$ ,  $B_{i,j}$  for  $i + j \leq 7$ , and  $N_{i,j,k}$  for  $i + j + k \leq 7$ . Among these,  $\Gamma_1$  was proved in [4],  $P_9$  was proved in [2], and the proof for  $N_{i,j,k}$  with i + j + k = 7 was completed in this series of two papers. The best known explicit result for  $Z_i$  is  $Z_3$  [6]; however, the proof for  $N_{1,1,5}$ in this paper gives implicitly  $Z_5$ , leaving open  $Z_i$  for  $6 \leq i \leq 7$ . Similarly, the best known explicit result for  $B_{i,j}$  is  $B_{1,2}$  (also in [6]); however, the results for  $N_{i,j,k}$  with i + j + k = 7 give implicitly all possible  $B_{i,j}$  with  $i + j \leq 6$ , leaving open  $B_{i,j}$  for i + j = 7. Finally,  $\Gamma_3$ , and  $\Gamma_5$ for  $n \geq 21$  remain open.

We summarize this discussion of possible graphs Y implying a 3-connected  $\{K_{1,3}, Y\}$ -free graph to be Hamilton-connected in the following table.

The graph $Y$	Possible	Known	Open
$\Gamma_i$	$\Gamma_1, \Gamma_3, \Gamma_5 \text{ for } n \geq 21$	$\Gamma_1$	$\Gamma_3$ ; $\Gamma_5$ for $n \ge 21$
$P_i$	$4 \le i \le 9$	$P_9$	
$Z_i$	$i \leq 7$	$Z_5$	$Z_6, Z_7 \text{ for } n \ge 21$
$B_{i,j}$	$i+j \leq 7$	$i+j \le 6$	i+j=7
$N_{i,j,k}$	$i+j+k \leq 7$	$i+j+k \le 7$	

**3.** Notice. During the reference process of this paper, sharp results for  $Z_i$  and  $B_{i,j}$  were proved in [18] and [19]. Namely, the following was shown (here  $W^+$  denotes the graph obtained from the Wagner graph W by attaching exactly one pendant edge to each of its vertices).

- [18] If G is a 3-connected  $\{K_{1,3}, Z_7\}$ -free graph that is not isomorphic to the graph  $L(W^+)$ , then G is Hamilton-connected. Specifically, every 3-connected  $\{K_{1,3}, Z_6\}$ -free graph is Hamilton-connected.
- [19] If G is a 3-connected  $\{K_{1,3}, B_{i,j}\}$ -free graph with  $i + j \leq 7$ , then G is Hamilton-connected.

Thus, the only remaining open cases are the graphs  $\Gamma_3$  (for all graphs), and  $\Gamma_5$  for  $|V(G)| \ge 21$ .

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